

## Using the POSIX API

Using the POSIX API
Threads, real-time and IPC



### The pthread library

- In Linux, when a new process is created, it already contains a thread, used to execute the main() function
- Additional threads can be created using the pthread library, which is part of the C library
- Of course all threads inside a given process will share the same address space, the same set of open files, etc.
- The pthread library also provide thread synchronization primitives: mutexes and conditions
- This pthread library has its own header: pthread.h
- Applications using pthread function calls should be explicitly linked with the pthread library gcc -o app app.c -lpthread



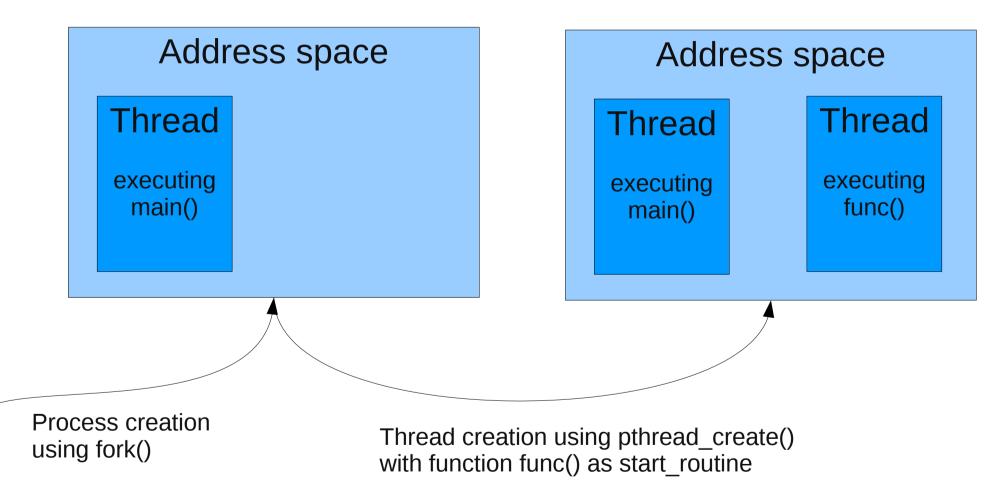
## Creating a new thread

► The function to create a new thread is pthread\_create()

- thread is a pointer to a pthread\_t structure that will be initialized by the function. Later, this structure can be used to reference the thread.
- Attr is a pointer to an optional structure pthread\_attr\_t. This structure can be manipulated using pthread\_attr\_\*() functions. It can be used to set various attributes of the threads (detach policy, scheduling policy, etc.)
- start\_routine is the function that will be executed by the thread
- arg is the private data passed as argument to the start\_routine function



# Creating a new thread (2)





#### Thread creation, code sample

```
#include <pthread.h>
void *thread(void *data)
  while(1) {
     printf(« Hello world from thread »);
int main(void) {
  pthread t th;
  pthread create(& th, NULL, thread, NULL);
  return 0;
```



#### Joinable and detached threads

- When the main() function exits, all threads of the application are destroyed
- ► The pthread\_join() function call can be used to suspend the execution of a thread until another thread terminates. This function must be called in order to release the ressources used by the thread, otherwise it remains as zombie.
- Threads can also be detached, in which case they become independent. This can be achieved using
  - Thread attributes at thread creation, using pthread\_attr\_setdetachstate(& attr, PTHREAD CREATE DETACHED);
  - pthread\_detach(), passing the pthread\_t structure as argument



#### Thread join, code sample

```
#include <pthread.h>
void *thread(void *data)
  int i;
  for (i = 0; i < 100; i++) {
     printf(« Hello world from thread »);
int main(void) {
  pthread t th;
  pthread create(& th, NULL, thread, NULL);
  pthread join(& th, NULL);
  return 0;
```



#### Thread cancelation

▶ It is also possible to cancel a thread from another thread using the pthread\_cancel() function, passing the pthread\_t structure of the thread to cancel.

```
#include <pthread.h>
void *thread(void *data)
   while(1) {
      printf(« Hello world from thread »);
}
int main(void) {
   pthread t th;
   pthread create(& th, NULL, thread, NULL);
   sleep(1);
   pthread cancel(& th);
   pthread join(& th, NULL);
   return 0;
```



## pthread mutexes (1)

- The pthread library provides a mutual exclusion primitive, the pthread\_mutex.
- Declaration and initialization of a pthread mutex
  - Solution 1, at definition time
    pthread\_mutex\_t lock = PTHREAD\_MUTEX\_INITIALIZER;
  - Solution 2, at runtime
     pthread\_mutex\_t lock;
     ...
     pthread\_mutex\_init(& lock, NULL);
     ...
     pthread mutex destroy(& lock);
  - ► The second argument to pthread\_mutex\_init() is a set of mutex-specific attributes, in the form of a pthread\_mutexattr\_t structure that can be initialized and manipulated using pthread mutexattr \*() functions.



### pthread mutexes (2)

- Take the mutex
  ret = pthread\_mutex\_lock(& lock);
- If the mutex is already taken by the calling threads, three possible behaviours depending on the mutex type (defined at creation time)
  - Normal (« fast ») mutex : the function doesn't return, deadlock
  - « Error checking » mutex : the function return with the -EDEADLK error
  - « Recursive mutex »: the function returns with success
- Palease the mutex
  ret = pthread\_mutex\_unlock(& lock);
- Try to take the mutex ret = pthread\_mutex\_trylock(& lock);



#### pthread conditions

- Conditions can be used to suspend a thread until a condition becomes true, as signaled by another thread.
- Initialization, static or dynamic
  - pthread\_cond\_t cond = PTHREAD\_COND\_INITIALIZER;
  - pthread\_cond\_t cond;
    pthread\_cond\_init(& cond, NULL);
- Wait for the condition pthread\_cond\_wait(& cond, & mutex) The mutex will be released before waiting and taken again after the wait
- Signaling the condition
  - To one thread waiting, pthread cond signal(& cond);
  - To all threads waiting, pthread cond broadcast(& cond);



#### pthread conditions example

# Receiver side

```
pthread_mutex_lock(& lock);
while(is_queue_empty())
   pthread_cond_wait(& cond, & lock);

/* Something in the queue,
   and we have the mutex ! */
pthread_mutex_unlock(& lock);
```

# Sender side

```
pthread_mutex_lock(& lock);

/* Add something to the queue */
pthread_mutex_unlock(& lock);

pthread_cond_signal(& cond);
```



#### Managing real-time priorities

See http://free-electrons.com/docs/realtime/ for an introduction C API Available through <sched.h> (see man sched.h for details)

- sched\_getscheduler, sched\_setscheduler
  Get / set the scheduling class of a process
- sched\_getparam, sched\_setparam
  Get / set the priority of a process
- sched\_get\_priority\_max, sched\_get\_priority\_min Get the maximum / minimum priorities allowed for a scheduling class.
- Sched\_rr\_get\_interval
  Get the current timeslice of the SCHED RR process
- sched\_yield
  Yield execution to another process.

Can also be manipulated from scripts with the chrt command.



### POSIX shared memory (1)

A great way to communicate between processes without going through expensive system calls.

- Dopen a shared memory object:
   shm\_fd = shm\_open("acme", O\_CREAT | O\_RDWR, 0666);
  A zero size /dev/shm/acme file appears.
- Set the shared memory object size ftruncate(shm\_fd, SHM\_SIZE); /dev/shm/acme is now listed with the specified size.
- If the object has already been sized by another process, you can get its size with the fstat function.



### POSIX shared memory (2)

- Map the shared memory in process address space: addr = mmap (0, SHM\_SIZE, PROT\_WRITE, MAP\_SHARED, shm\_fd, 0); Now we have a memory area we can use!
- Lock the shared memory in RAM (best for real-time tasks): mlock(addr, SHM SIZE);
- Use the shared memory object! Other processes can use it too.



### POSIX shared memory (3)

#### **Exiting**

Unmap the shared memory object: munmap (addr, SHM\_SIZE); This automatically unlocks it too.

```
Close it:
  close (shm_fd);
```

Remove the shared memory object:

```
shm_unlink ("acme");
```

The object is effectively deleted after the last call to shm\_unlink.

More details in man shm\_open.



### POSIX message queues

Deterministic and efficient IPC. See man mqueue.h. Advantages for real-time applications:

- Preallocated message buffers
- Messages with priority.
  A message with a higher priority is always received first.
- Send and receive functions are synchronous by default. Possibility to set a wait timeout to avoid non-determinism.
- Support asynchronous delivery notifications.



#### Creating and opening a message queue

Declare queue attributes:

```
queue_attr.mq_maxmsg = 16;
  /* max number of messages in queue */
queue_attr.mq_msgsize = 128;
  /* max message size */
```

Open a queue:



#### Posting a message

Posting a message:
 #define PRIORITY 3
 char msg[] = "Goodbye Bill";

```
mqsend(qd, msg, strlen(msg), PRIORITY);
```

Closing the queue: mq close(qd);

Caution: simplistic example code. Should check return values.



#### Receiving a message

#### From another application:

Opening the shared message queue:

Waiting for a message:

```
mq_receive(qd, text, buf, buf_size, &prio);
```

Close the queue:

```
mq close(qd);
```

Destroy the queue:

```
mq unlink("/msg queue");
```



## POSIX semaphores (1)

Resources for sharing resources between threads or processes. See man semaphore.h.

- Named semaphores: can be used between unrelated processes.
- ▶ Unnamed semaphores: can be used between threads from the same process, or by related processes (parent / child).



# POSIX semaphores (2)

- Sem\_open
  Open and / or create
  a named semaphore.
- sem\_close
  Close a named semaphore
- sem\_unlink
  Destroy a named semaphore
- sem\_init
  Initialize an unnamed semaphore
- b sem\_destroy
  Destroy an unnamed semaphore

- sem\_getvalue
  Get current semaphore count
- Try to lock the semaphore.
  Wait otherwise.
- Just tries to lock the semaphore, but gives up if the semaphore is already locked.
- sem\_post
  Release the semaphore.



#### POSIX signals

- Signals are a mechanism to notify a process that an event occured: expiration of a timer, completion of an asynchronous I/O operation, or any kind of event specific to your application
- Signals are also used internally by the system to tell a process that it must be suspended, restarted, stopped, that is has done an invalid memory reference, etc.
- Each signal is identified by a number: SIGSEGV, SIGKILL, SIGUSR1, etc.
- An API is available to catch signals, wait for signals, mask signals, etc.
- See signal(7) for a general description of the signal mechanism



# Registering a signal handler

- A signal handler can be registered using
  - > sighandler\_t signal(int signum, sighandler\_t
    handler);
    - The handler has the following prototype: void handler(int signum)
  - int sigaction(int signum, const struct sigaction
    \*act, struct sigaction \*oldact);
    - The sigaction structure contains the reference to the handler
    - The handler can have two different prototypes
      - void handler(int signum)
      - void handler(int signum, siginfo\_t \*info, void \*data)
- Inside the handler code, only some functions can be used: only the async-signal-safe functions, as documented by signal(7).



## Signal registration example

```
#include <signal.h>
#include <assert.h>
#include <unistd.h>
#include <stdio.h>
void myhandler(int signum)
{
    printf("Signal catched!\n");
}
int main(void)
{
    int ret;
    struct sigaction action = {
        .sa handler = myhandler,
    };
    ret = sigaction(SIGUSR1, & action, NULL);
    assert(ret == 0);
    while(1);
    return 0;
}
```

From the command line, the signal can then be sent using kill -USR1 PID



#### Sending a signal

- From the command line, with the famous kill command, specifying the PID of the process to which the signal should be sent
  - By default, kill will send SIGTERM
  - Another signal can be sent using kill -USR1
- POSIX provides a function to send a signal to a process
  - int kill(pid\_t pid, int sig);
  - In a multithread program, the signal will be delivered to an arbitrary thread. Use tkill() to send the signal to a specific thread.



#### Signal sets and their usage

- A type sigset\_t is defined by POSIX, to hold a set of signals
- This type is manipulated through different functions
  - sigemptyset() to empty the set of signals
  - sigaddset() to add a signal to a set
  - sigdelset() to remove a signal from a set
  - sigfillset() to fill the set of signals with all signals
- Signals can then be blocked or unblocked using sigprocmask(int how, const sigset\_t \*set, sigset\_t \*oldset);
- sigset\_t are also used in many other functions
  - sigaction() to give the list of signals that must be blocked during execution of the handler
  - sigpending() to get the list of pending signals



### Waiting for signals

#### 2 ways of waiting for signals:

- sigwaitinfo() and sigtimedwait() to wait for blocked signals (signals which remain pending until they are processed by a thread waiting for them.)
- sigsuspend() to register a signal handler and suspend the thread until the delivery of an unblocked signal (which are delivered without waiting for a thread to wait for them).



## POSIX real-time signals

#### Regular signals

- Just 2 applications-specific signals: SIGUSR1 and SIGUSR2
- No signal priorities
- Signals can't carry any extra information.
- Signals can be lost. When a signal is sent multiple times, the receiver will just process one instance.

#### POSIX signals

- Whole range of application specific signals: SIGRTMIN to SIGRTMAX
- Priorities available.Top priority signals delivered first.
- Possible to carry extra information in a signal.
- Signals are queued. All pending signals are processed: no signal is lost.



#### POSIX clocks and timers

#### Compared to standard (BSD) timers in Linux

- Possibility to have more than 1 timer per process.
- Increased precision, up to nanosecond accuracy
- Timer expiration can be notified either with a signal or with a thread.
- Several clocks available.



#### Available POSIX clocks (1)

#### Defined in /usr/include/linux/time.h

- CLOCK\_REALTIME System-wide clock measuring the time in seconds and nanoseconds since Jan 1, 1970, 00:00. Can be modified. Accuracy: 1/HZ (1 to 10 ms)
- ► CLOCK\_MONOTONIC

  System-wide clock measuring the time in seconds and nanoseconds since system boot. Cannot be modified, so can be used for accurate time measurement.

  Accuracy: 1/HZ



#### Available POSIX clocks (2)

- CLOCK\_PROCESS\_CPUTIME\_ID
  Measures process uptime. 1/HZ accuracy. Can be changed.
- CLOCK\_THREAD\_CPUTIME\_ID Same, but only for the current thread.



### Time management

#### Functions defined in time.h

- clock\_settime
  Set the specified clock to a value
- clock\_gettime
  Read the value of a given clock
- clock\_getres
  Get the resolution of a given clock.

See man time.h and the manual of each of these functions.



#### Using timers (1)

#### Functions also defined in time.h

- Clock\_nanosleep Suspend the current thread for the specified time, using a specified clock.
- nanosleep Same as clock\_nanosleep, using the CLOCK REALTIME clock.



#### Using timers (2)

- timer\_create
  Create a timer based on a given clock.
- timer\_delete
  Delete a timer
- timer\_settime
  Arm a timer.
- timer\_gettime
  Access the current value of a timer.



#### Using high resolution timers

- Available in Linux since 2.6.21 (on x86).
  Now available on most supported platforms.
- ▶ Depending on the hardware capabilities, this feature gives microsecond or nanosecond accuracy to the regular clocks (CLOCK\_REALTIME, CLOCK\_MONOTONIC).
- No need to recompile your applications!



#### Asynchronous I/O

- Helpful to implement non-blocking I/O.
- Allows to overlap compute tasks with I/O processing, to increase determinism.
- Supported functionality:
  - Send multiple I/O requests at once from different sources
  - Cancel ongoing I/O requests
  - Wait for request completion
  - Inquire the status of a request: completed, failed, or in progress.
- API available in aio.h (man aio.h for details)



#### Compiling instructions

- Includes: nothing special to do. Available in the standard path.
- Libraries: link with librt
- Example:
   gcc -lrt -o rttest rttest.c



#### POSIX manual pages

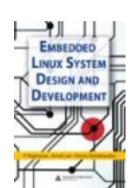
POSIX manual pages may not be installed on your system

- On Debian Linux, based systems, to find the names of the corresponding packages: apt-cache search posix Then, install these packages as follows: apt-get install manpages-posix manpages-posix-dev
- Other distributions should have similar package names.
- These manual pages are also available on-line: http://www.opengroup.org/onlinepubs/009695399/idx/realtime.html

You can almost consider these manual pages as specifications. The standard can also be accessed on <a href="http://www.unix.org/online.html">http://www.unix.org/online.html</a> (registration required).



#### More information on the POSIX interface



- The POSIX manual pages
- Embedded Linux System Design and Development P. Raghavan, A. Lad, S. Neelakandan, Auerbach, Dec. 2005. http://free-electrons.com/redirect/elsdd-book.html Very nice and clear coverage on real-time programming with the POSIX interface. Nice and useful examples.
- Guide to real-time programming http://www.phys.uu.nl/DU/unix/HTML/APS33DTE/TITLE.HTM A 11-year old document, with some Digital Unix specifics, but still up to date (thanks to standards).